#### Lecture 9 - Discrete MDPs

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# Recap: Discrete-time Markov Decision Process (MDP)

Discrete-time Markov decision process (MDP), denoted as the tuple  $(S, A, T, R, \rho_0, \gamma)$ .

- $\mathcal{S}$  the state space;
- $\mathcal{A}$  the action space.  $\mathcal{A}$  can depend on the state s for  $s \in \mathcal{S}$ ;
- $\mathcal{T}: \mathcal{S} \times \mathcal{A} \rightarrow \Delta(\mathcal{S})$  the environment transition probability function;
- $\mathcal{R}: \mathcal{S} \times \mathcal{A} \to \Delta(\mathbb{R})$  the reward function;
- $ho_0 \in \Delta(\mathcal{S})$  the initial state distribution;
- $\gamma \in [0,1]$  the discount factor.

Note that  $\Delta(\mathcal{X})$  denotes the set of all distributions over set  $\mathcal{X}_{max}$ 



## Recap: Discrete-time Markov Decision Process (MDP)

A stationary MDP follows for t = 0, 1, ... as below, starting with  $s_0 \sim \rho_0$ .

- The agent observes the current state *s<sub>t</sub>*;
- The agent chooses an action  $a_t \sim \pi(a_t \mid s_t)$ ;
- The agent receives the reward  $r_t \sim \mathcal{R}(s_t, a_t)$ ;
- The environment transitions to a subsequent state according to the Markovian dynamics s<sub>t+1</sub> ~ T(s<sub>t</sub>, a<sub>t</sub>).

This process generates the sequence  $s_0, a_0, r_0, s_1, \ldots$  indefinitely. The sequence up to time *t* is defined as the trajectory indexed by *t*, as  $\tau_t = (s_0, a_0, r_0, s_1, \ldots, r_t)$ .  $\stackrel{\text{$\clubsuit$}}{\Rightarrow} \psi \gtrsim \chi \not\in (\mathcal{R} \hspace{1pt} \mathfrak{I})$ The Chinese University of Hong Kong, Shenzhen

## Recap: Discrete-time Markov Decision Process (MDP)

The goal is to optimize the expected discounted cumulative return

$$\mathbb{E}_{s_t,a_t,r_t,t\geq 0}\left[R_0\right] = \mathbb{E}_{s_t,a_t,r_t,t\geq 0}\left[\sum_{t=0}^{\infty} \gamma^t r_t\right]$$

over the agent's policy  $\pi$ .



A Markov chain, also known as a homogeneous Markov chain, refers to an infinite process  $x_1, \ldots, x_T, \ldots$  where

$$\mathbb{P}(x_{t+1} \mid x_t, \ldots, x_1) = \mathbb{P}(x_{t+1} \mid x_t) = \mathbb{P}_{\mathcal{M}}(x' \mid x)$$

holds almost surely for some probability measure  $\mathbb{P}_{\mathcal{M}}$ .



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Some key definitions:

- The state space S is countable, e.g., state space [n] = [1, 2, ..., n] is finite.
- The transition probability matrix is  $P \in \mathbb{R}^{n \times n}$  where the element  $P_{ii'}$  on the *i*-th and *i'*-th column equals  $\mathbb{P}(x' = i' \mid x = i)$ . It is  $P^{\pi}$  in MDPs.
- The state value function is V(s), and the value vector is  $V \in \mathbb{R}^n$ .
- The reward function is  $\mathcal{R}(s)$ , and the reward vector is  $r \in \mathbb{R}^n$ .
- The initial state distribution  $ho_0 \in \mathbb{R}^n$ .

When  $\mathcal{R}(s)$  is deterministic, r is a deterministic vector. The Bellman equation then writes  $V = r + \gamma P V$ . Since P is a Markov matrix,  $I - \gamma P$  is invertible when  $\gamma < 1$  and the value function can be solved by

 $V=(I-\gamma P)^{-1}r.$ 



**Reducible states.** For two states  $i, i' \in [n]$ , if there exists a T such that  $\mathbb{P}(i' \in \{s_1, \ldots, s_T\} \mid s_0 = i) > 0$ , we say that i' is accessible from i. If i is accessible from i' and i' is accessible from i, we say that i and i' communicate with each other. If for any  $i, i' \in [n]$ , i and i' communicate with each other, the Markov chain is irreducible.

• For a Markov chain that is reducible, it is intuitive to partition the chain into irreducible components (likewise, to consider each connected component in a graph). It is therefore sensible to assume that the Markov chain is irreducible.



**Periodicity.** For  $i \in [n]$ , the period of state *i* is the largest integer *d* satisfying  $\mathbb{P}(s_t \neq i \mid s_0 = i, t \neq 0 \mod d)$ , or infinity if such a largest integer does not exist. When d = 1, state *i* is aperiodic, and otherwise, state *i* is periodic with period *d*.

- In an irreducible Markov chain, all states have the same period. An irreducible Markov chain is aperiodic if all states are aperiodic.
- Mathematically, a chain is aperiodic if and only if  $P^t$  contains only positive elements for some positive integer t.



Ergodicity. A Markov chain that is irreducible and aperiodic must be ergodic. We commonly assume a chain to be ergodic without loss of generality.

• For the rest of the course, unless otherwise specified, we assume the Markov chains to be ergodic. In MDPs however, in general, there exist policies such that the chain induced by the policies are not ergodic.



# Policy Evaluation With a Known Model

#### Model-based Policy Evaluation:

- When at least one of *P* and *r* is known, the problem is policy evaluation with a known model.
- When both P and r are unknown we can make an effort to estimate a P' such that P and P' are close in some measure of discrepancy (or r', respectively).
  If otherwise and we only utilize the access to the environment transition, the method is categorized as model-free policy evaluation.



## Policy Evaluation With a Known Model

**Policy Evaluation (PE)**: compute the value function given a fixed policy. When at least one of P and r is known, the problem is policy evaluation with a known model.

- Under discrete state and action spaces,
- When both *P* and *r* are known.
- When  $\gamma < 1$

The solution is:  $V = (I - \gamma P)^{-1}r$ .



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## Model-based Policy Evaluation

When both P and r are unknown and we estimate a P' such that P and P' are close (or r', respectively), the problem is model-based policy evaluation.

#### Lemma

Assume that  $0 \le r \le 1$ . Let  $\varepsilon \in (0, \frac{1}{1-\gamma})$ . There is an absolute constant c such that once one has collected at least:

$$N \geq rac{\gamma}{(1-\gamma)^4} rac{n^2 m \log(cnm/\delta)}{arepsilon^2}$$

Some assumptions:

- Let the reward function be deterministic.
- Let the reward is bounded by [0,1] in discrete MDPs.
- Let the reward function  $\mathcal{R}(s, a)$  defined by a matrix  $r \in \mathbb{R}^{n \times m}$ , where the element at the *i*-th row and the *j*-th column denotes  $\mathcal{R}(i, j)$ .
- Let  $P_j$  be the transition matrix for the policy that chooses action j at every state. Recall that in discrete MDPs a value function V is optimal if and only if the Bellman optimality equation:

$$V = \max_{j} [r_j + \gamma P_j V]$$

is satisfied with  $P \in \{P_1, \ldots, P_m\}$ .



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By exhausting the action set under the max operator and numbering the actions from 1 to m, the Bellman optimality equation is formulated into the below linear program:

minimize 
$$e^T V$$
  
subject to  $(I - \gamma P_j)V - r_j \ge 0$ ,  $j = 1, ..., m$ ,

where e is the all-one vector and  $e^T V$  is a dummy objective. Linear programming is in P and can be solved in poly(n, m). We consider a problem solved if we can cast it to a linear program. Though, this requires  $P_i$  to be known.



The dual of the above linear program is

$$\begin{array}{ll} \underset{\lambda_{1},\ldots,\lambda_{m}}{\operatorname{maximize}} & \sum_{j} \lambda_{j}^{T} r_{j} \\ \text{subject to} & \sum_{j} (I - \gamma P_{j}^{T}) \lambda_{j} = \mathsf{e}, \\ & \lambda_{j} \geq \mathsf{0}, \quad j = 1, \ldots, m \end{array}$$

The dual formulation optimizes  $\lambda_1, \ldots, \lambda_m$ , which could represent policy information (More specifically, occupancy measure:  $P(s_t = s, a_t = a | \mu_0, \pi, \mathcal{T})$ ). 香港中文大學(深圳) The Chinese University of Hong Kong, Shenzhen

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#### Lemma

There exists an optimal dual solution  $\lambda_j^*$ , j = 1, ..., m, an optimal deterministic policy  $\pi^*(\cdot)$ , and the corresponding transition matrix  $P^*$ , such that

$$\sum_j \lambda_j^* = (I - \gamma P^{*T})^{-1}$$
e,

and the *i*-th entry of  $\lambda_j^*$  equals to the *i*-th entry of  $\sum_j \lambda_j^*$  if  $\pi^*(i) = j$ , and zero otherwise.



Lemma

The  $\ell^1$ -norm  $\|\sum_j \lambda_j^*\|_1$  of the dual optimum is exactly  $n/(1-\gamma)$ .



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#### Lemma

The stochastic policy  $\pi(j \mid i) = \lambda_j^{\prime(i)} / \sum_{j'} \lambda_{j'}^{\prime(i)}$  achieves a value V' such that  $e^T V' = \sum_j \lambda_j^{\prime T} r_j$ .



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## Question and Answering (Q&A)



